

Neighbour-transitive codes in generalised quadrangles

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Combinatorial Designs and Codes, Rijeka (virtually)

Codes in graphs

Delsarte and Biggs independently introduced the concept of a code in a graph in 1973

- graph Γ with vertex set $V\Gamma$
- code C in Γ is a subset of $V\Gamma$
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Graphs most commonly considered are the Johnson and Hamming graphs

Codes in graphs

Let C be a code in a graph Γ

Important parameters:

- *minimum distance* of C : δ
smallest distance between distinct codewords
- *covering radius* of C : ρ
largest distance between a vertex of Γ and its nearest codeword

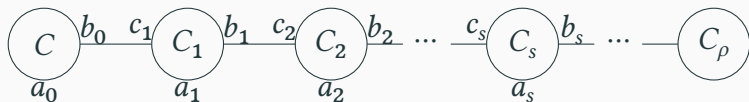
s -Regular codes

A code C in Γ gives rise to a partition $\{C_i\}_{i=0}^{\rho}$ of $V\Gamma$, where C_i is the set of vertices distance i to their nearest codeword.



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Definition (Delsarte 1973)

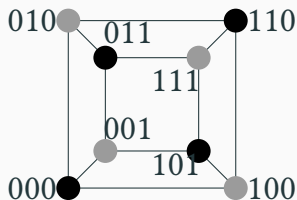
A code C is s -regular if for all $i \leq s$ there exist integers a_i, b_i, c_i (as above) depending only on i giving the number of adjacent vertices for each $\alpha \in C_i$.

A code that is ρ -regular is called *completely regular*.

Automorphisms Groups

Definition

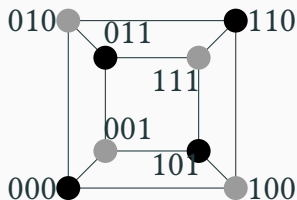
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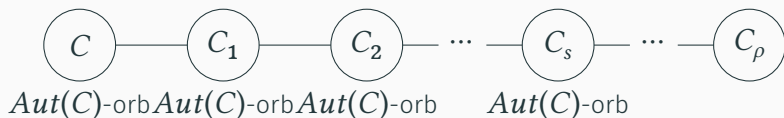
E.g. Interchanging the inner and outer pairs of vertices is an automorphism of $H(3, 2)$ but not of C , whilst a reflection on a diagonal is an automorphism of both.

s -Neighbour-transitive codes

Definition

A code C is s -neighbour-transitive (s -NT) if $\text{Aut}(C)$ acts transitively on C_i for each $i \leq s$.

If C is ρ -NT then C is said to be *completely transitive* (CT).

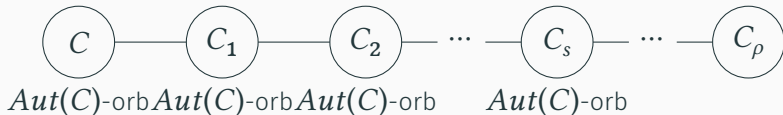


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CT-ity originally introduced by Solé (1987) – above definition due to Giudici and Praeger (2000).

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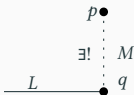
Work has been done on:

- NT codes in Johnson graphs – Liebler and Praeger, Neunhoffer and Praeger, Iopollo's PhD thesis.
- CT codes in Hamming graphs – Giudici and Praeger, Gillespie's PhD thesis, Bailey and Hawtin.
- NT codes in Hamming graphs – Gillespie's PhD thesis.
- 2-NT codes in Hamming graphs – H. thesis; Gillespie, Giudici, H. and Praeger.

Generalised quadrangles

A *generalised quadrangle* (shortly a GQ) – introduced by Tits (1959) – is an incidence structure $\mathcal{Q} = (\mathcal{P}, \mathcal{L}, \mathbf{I})$, where \mathcal{P} and \mathcal{L} are sets called points and lines, respectively, and \mathbf{I} is a symmetric point-line incidence relation such that:

1. Each point is incident with $t + 1$ lines ($t \geq 1$) and two distinct points are incident with at most one line.
2. Each line is incident with $s + 1$ points ($s \geq 1$) and two distinct lines are incident with at most one point.
3. If p is a point and L is a line not incident with p , then there is a unique pair $(q, M) \in \mathcal{P} \times \mathcal{L}$ for which $p \mathbf{I} M \mathbf{I} q \mathbf{I} L$.



Classical GQs

Classical GQs are associated with certain classical groups and isotropic points and lines of projective geometries under certain bilinear/sesquilinear forms.

\mathcal{Q}	order	\mathcal{Q}^D	$\text{Aut}(\mathcal{Q})$
$Q_4(q)$	(q, q)	$W_3(q)$	$P\Gamma O_5(q)$
$Q_5^-(q)$	(q, q^2)	$H_3(q^2)$	$P\Gamma O_6^-(q)$
$W_3(q)$	(q, q)	$Q_4(q)$	$P\Gamma Sp_4(q)$
$H_3(q^2)$	(q^2, q)	$Q_5^-(q)$	$P\Gamma U_4(q)$
$H_4(q^2)$	(q^2, q^3)	$H_4(q^2)^D$	$P\Gamma U_5(q)$

We will be interested mainly in $W_3(q)$.

Note that there are various other constructions for GQs, but we won't consider these.

Point-line incidence graph of a GQ

The point-line incidence graph Γ of a GQ $\mathcal{Q} = (\mathcal{P}, \mathcal{L}, \mathbf{I})$ is defined on the vertex set $V\Gamma = \mathcal{P} \cup \mathcal{L}$ with adjacency given by \mathbf{I} .

Properties of Γ :

- bipartite
- diameter 4
- girth 8

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An code C in Γ with $\delta = 4$ is either a *partial ovoid* or a *partial spread*, if $\rho = 2$ then C is either an *ovoid* or a *spread*, and if $\rho = 3$ then C is either a *maximal partial ovoid* or a *maximal partial spread*.

Ovoids and spreads

Theorem (Crnković , H., Švob)

Let C be a NT code with $\delta = 4$ and $\rho = 2$ in a thick classical GQ and assume that $\text{Aut}(C)$ is insoluble. Then C is equivalent to one of the following:

1. The regular spread of $W_3(q)$.
2. A classical ovoid of $H_3(q^2)$.

Proof applies a theorem of Bamberg and Pentilla on transitive m -systems in polar spaces.

Examples in $W_3(q)$

- $V \cong \mathbb{F}_q^4$ equipped with the symplectic form
 $f(x, y) = x_1y_2 - x_2y_1 - x_3y_4 + x_4y_3$
- G a sharply transitive subgroup of $SL_2(q)$ (below)
- $C = \{ \begin{bmatrix} I & A \end{bmatrix} \mid A \in G \}$
- note: $\det A = 1$ means form gives 0 and so the rowspaces of elements of C represent lines of $W_3(q)$

q	G
2	$GL_1(4)$
3	Q_8
5	$2.A_4$
7	$2.S_4$
11	$SL_2(5)$

These maximal partial spreads of $W_3(q)$ were originally found by Penttila.

Main result

Theorem (Crnković , H., Švob)

Let C be a NT code with minimum distance 4 in the generalised quadrangle $W_3(q)$. Then the following hold:

1. C is a regular spread $\iff \rho = 2$.
2. $|C| = q^2$ implies $\rho = 4$ and C is a spread minus a line.
3. $|C| = q^2 - 1$ implies $q = 2, 3, 5, 7$ or 11 and C is one of the codes from the previous slide, with $\rho = 3$.
4. $|C| = q + 1$ and $\rho = 3$ implies C is the set of points on a hyperbolic line.
5. If $q = 3$ and $|C| = 5$ then there is a unique code with $\rho = 3$.

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Conjecture (Crnković , H., Švob)

Let C be a NT maximal partial ovoid or maximal partial spread of $W_3(q)$. Then C is as in parts 1,3,4 or 5 above.

Thanks!

Thanks for listening!