

# **Pless symmetry codes, ternary QR codes, and related Hadamard matrices and designs**

Vladimir D. Tonchev

Michigan Technological University



**Vera Pless (1931–2020)**

**ON A NEW FAMILY OF SYMMETRY CODES AND  
RELATED NEW FIVE-DESIGNS**

BY VERA PLESS

Communicated by Wallace Givens, May 15, 1969

For every prime  $p \equiv -1 \pmod{3}$  we define a self-orthogonal  $(2p+2, p+1)$  code over GF(3). It can be shown that the group leaving a  $(2p+2, p+1)$  code invariant is  $PSL_2(p)$ . The minimum weights of the first five codes in the family are determined and lead to new 5-designs.

Let  $t, r,$  and  $n$  be integers with  $t \leq r \leq n$ . A  $\lambda; t-r-n$  design  $D$  is a collection of subsets of the  $n$  integers, each subset containing  $r$  elements, such that any  $t$ -subset of the  $n$  integers is contained in the same number  $\lambda$  of subsets in  $D$ . Some designs, a  $1; 5-6-12$ , a  $1; 5-8-24$ , and a  $48; 5-12-24$  associated with the Mathieu groups  $M_{12}$  and  $M_{24}$ , have been known for a long time. Recently, [1] and [5],  $2; 5-6-12$  and  $2; 5-8-24$  designs have been found. Using coding theory [2] other 5-designs were found for  $n=24$  and  $n=48$ . We have found new 5-designs for  $n=36$  and  $n=60$  and a number of  $r$ 's. Also we found new 5-designs for  $n=24$  and  $n=48$  which are not equivalent to the ones mentioned above. Two  $t$ -designs are called equivalent if there is a permutation of the  $n$  integers so that the subsets of  $D$  go onto subsets in  $D$ .

Let  $V_{2p+2}$  be a vector space over GF(3) with a fixed, orthonormal basis. We call a subspace of this space an error correcting code. We define a family of codes of  $\dim(p+1)$  (referred to as  $(2p+2, p+1)$  codes) by a basis  $(I, S_p)$  where  $S_p$  is given below.

$$S_p = \begin{array}{c} \infty \\ 0 \\ 1 \\ \vdots \\ i \\ \vdots \\ (p-1) \end{array} \left| \begin{array}{cccccc} \infty & 0 & 1 & \cdots & j & \cdots & (p-1) \\ 0 & 1 & 1 & 1 & 1 & 1 & \\ \chi(-1) & \chi(0) & \chi(1) & \chi(j) & \chi(p-1) & & \\ \chi(-1) & & & & & & \\ \chi(-1) & & & \chi(j-i) & & & \\ \chi(-1) & & & & & & \end{array} \right.$$

where  $\chi(0) = 0, \chi(\text{a square}) = 1, \chi(\text{a nonsquare}) = -1$ . We refer to the code generated by  $(I, S_p)$  as  $C(p)$ .

## Symmetry Codes over GF(3) and New Five-Designs

VERA PLESS

*Air Force Cambridge Research Laboratories, L. G. Hanscom Field,  
Bedford, Massachusetts*

*Communicated by Andrew Gleason*

Received November 3, 1969

For every odd prime power  $q$  where  $q \equiv -1(3)$  we define a  $(2q + 2, q + 1)$  code over the field of three elements. It is shown that all the codes in this family are self orthogonal.

For  $q = 5$ , the  $(12, 6)$  code is equivalent to the extended Golay code. For  $q = 11$ , it can be shown that the minimum weight of the  $(24, 12)$  code is 9. For  $q = 17, 23, 29$  it is shown, in part by computer, that the minimum weights of the  $(36, 18)$ ,  $(48, 24)$ , and  $(60, 30)$  codes are 12, 15, and 18 respectively.

There are 5-designs associated with vectors of certain weights in the  $(12, 6)$ ,  $(24, 12)$ ,  $(36, 18)$ ,  $(48, 24)$ , and  $(60, 30)$  codes. There are new 5-designs associated with the last four codes mentioned. The 5-designs related to the  $(36, 18)$  and  $(60, 30)$  codes are the first 5-designs found with their parameters.

For each  $q$  we construct a group  $P$  of  $(2q + 2) \times (2q + 2)$  monomial matrices. We show that  $P$  leaves the  $(2q + 2, q + 1)$  code in the family invariant, and that  $P/\{I, -I\}$  is isomorphic to  $\text{PGL}_2(q)$ .

We can form a Hadamard matrix by considering the rows of this matrix as certain maximal weight vectors contained in this code. This Hadamard matrix is left invariant by the group  $P$  described above.

### I. INTRODUCTION

In this paper we define a family of codes over GF(3), where each code is associated with  $q$ , a power of an odd prime, such that  $q \equiv -1(3)$ . The first code in the family is the well known Golay  $(12, 6)$  code. The next four codes have high minimum weights and new 5-designs are associated with them. We also describe a group which leaves each code invariant.

In Section II we define each code in terms of a basis  $[I, S_q]$  where  $S_q$  is given in terms of the residues and non-residues in GF( $q$ ). The matrix  $S_q$  figures prominently [5, pp. 209, 210] in the construction of Hadamard matrices. In Section II these codes are shown to be self orthogonal. Also it is shown that, for  $q \equiv 1(4)$ ,  $[-S_q, I]$  is also a basis of its code, and, for  $q \equiv 3(4)$ ,  $[S_q, I]$  is again a basis of its code.

In Section III, using this and other properties, we determine the

## New 5-Designs\*

E. F. ASSMUS, JR.<sup>†</sup> AND H. F. MATTSON, JR.

*Applied Research Laboratory, Sylvania Electronic Systems,  
Waltham, Massachusetts 02154*

*Communicated by A. M. Gleason*

Received April 4, 1968

### ABSTRACT

A  $t$ -design on a point-set  $S$  is a collection  $\mathcal{D}$  of subsets of  $S$ , all of the same cardinality, with the property that every  $t$ -subset of  $S$  is contained in precisely  $\lambda$  elements of  $\mathcal{D}$ ,  $\lambda$  a fixed integer parameter of the design. Via the theory of error-correcting codes, we construct here several new 5-designs on 24 and 48 points as well as the two classic 5-designs on 12 and 24 points associated with the Mathieu groups  $M_{12}$  and  $M_{24}$ . We are able, in many cases, to say what the automorphism groups of the new 5-designs are.

### 1. INTRODUCTION

Tactical configurations and Hadamard matrices, studied for many years by combinatorialists, and the newer subject called error-correcting codes, studied for less than twenty years, have some interesting interconnections. The purpose of this report is to establish a number of new results arising therefrom.

Our main result is the construction (via Theorem 4.2) of several new 5-designs on 24 and 48 points and the determination (Section 5) of their automorphism groups as  $\text{PSL}_2(23)$  and  $\text{PSL}_2(47)$ , respectively. A secondary result (Section 5) is that  $\text{PSL}_2(l)$  is the automorphism group of certain quadratic-residue codes of length  $l + 1$  for all primes  $l$  having  $(l - 1)/2$  prime and satisfying  $23 \leq l \leq 4,079$ . (For  $l = 23$  we use [15] and a new 5-design on 24 points; the other cases are an immediate consequence of the Parker and Nikolai search [22].) We have derived elsewhere [7] the consequence that for  $l \equiv -1 \pmod{12}$ , the Paley-Hadamard matrix of order  $l + 1$  has  $\text{PSL}_2(l)$  as automorphism group for  $l$  as above.

---

\* The research reported in this paper was sponsored by the Air Force Cambridge Research Laboratories, Office of Aerospace Research, under contract AF19(628)-5998.

<sup>†</sup> Lehigh University and Sylvania.

Let  $A$  and  $B$  be linear orthogonal  $(n, k)$  and  $(n, n - k)$  codes over  $\text{GF}(q)$  with minimum weights  $d$  and  $e$ . Let  $t$  be an integer less than  $d$ . Let  $v_0$  be the largest integer satisfying

$$v_0 - \left\lfloor \frac{v_0 + (q - 2)}{q - 1} \right\rfloor < d,$$

and  $w_0$  the largest integer satisfying

$$w_0 - \left\lfloor \frac{w_0 + (q - 2)}{q - 1} \right\rfloor < e,$$

where, if  $q = 2$ , we take  $v_0 = w_0 = n$ . Then two vectors of  $A$  with weight at most  $v_0$  having their non-0 coordinates in the same places must be scalar multiples of each other, and the same for  $B$ . This property is essential to our method of proof of our main result,

**THEOREM 4.2.** *Suppose that the number of non-0 weights of  $B$  which are less than or equal to  $n - t$  is itself less than or equal to  $d - t$ . Then, for each weight  $v$  with  $d \leq v \leq v_0$ , the vectors of weight  $v$  in  $A$  yield a  $t$ -design, and for each weight  $w$  with  $e \leq w \leq \min\{n - t, w_0\}$ , the vectors of weight  $w$  in  $B$  yield a  $t$ -design.*

Before proving the above result we remark that for  $B$  we will in fact show that for each weight  $w$ , with  $e \leq w \leq \min\{n - t, w_0\}$ , the vectors of weight  $w$  yield blocks the complements of which form a  $t$ -design. We will need the following combinatorial

**LEMMA.** *Suppose  $(S, \mathcal{D})$  is a  $t$ -design. Then, if  $T$  and  $T'$  are two  $t$ -subsets of  $S$ , and  $k$  an integer satisfying  $0 \leq k \leq t$ , we see that*

$$|\{D \in \mathcal{D}; |D \cap T| = k\}| = |\{D \in \mathcal{D}; |D \cap T'| = k\}|.$$

*That is, the number of subsets in  $\mathcal{D}$  intersecting a given  $t$ -subset in precisely  $k$  points is independent of the chosen  $t$ -subset.*

**PROOF:** For  $k = t$  the assertion is simply the condition that  $(S, \mathcal{D})$  is a  $t$ -design. Now we use induction downward observing that for  $K \subseteq T$ ,  $|K| = k$ , we see that

$$|\{D \in \mathcal{D}; K \subseteq D\}| = \frac{\lambda \binom{n-k}{t-k}}{\binom{d-k}{t-k}} = \lambda_k,$$

# Summary

- A code  $L(q)$ , monomially equivalent to the Pless symmetry code  $C(q)$  of length  $2q + 2$ , contains the  $(0,1)$ -incidence matrix of a **Hadamard 3- $(2q + 2, q + 1, (q - 1)/2)$  design**  $D(q)$  associated with a **Paley-Hadamard matrix of type II**.
- The ternary extended QR code of length  $n \equiv 0 \pmod{12}$  contains a **Hadamard 3-design** associated with a **Paley-Hadamard matrix of type I**.
- If  $q = 5, 11, 17, 23$ , the full **permutation** automorphism group of  $L(q)$  coincides with the **full** automorphism group of  $D(q)$ .
- A similar result holds for the ternary extended QR codes of lengths 24 and 48.
- All Hadamard matrices of order 36 formed by codewords of the Pless symmetry code  $C(17)$  are classified up to equivalence:
  - (1) the **Paley-Hadamard matrix  $H$  of type II**, with a full automorphism group of order 19584;
  - (2) a **regular Hadamard matrix  $H'$**  such that the related symmetric  $2$ - $(36, 15, 6)$  design has **trivial** automorphism group.

# MacWilliams Identity and Gleason's Theorem

## MacWilliams Identity

If  $A(x) = \sum_{i=0}^n A_i x^i$  and  $B(x) = \sum_{i=0}^n B_i x^i$  are the weight enumerators of a linear  $[n, k]_q$  code  $C$  and its dual code  $C^\perp$ , then

$$q^k B(x) = (1 + (q - 1)x)^n A\left(\frac{1 - x}{1 + (q - 1)x}\right).$$

## An upper bound (Mallows and Sloane 1973)

If  $C$  is a self-dual  $[n, n/2, d]$  ternary code then

$$d \leq 3\left[\frac{n}{12}\right] + 3.$$

## Definition

A ternary self-dual code of length  $n$  is **extremal** if it meets the upper bound:  $d = 3\left[\frac{n}{12}\right] + 3$ .



# The Assmus-Mattson Theorem

## Theorem (Assmus and Mattson 1969)

If  $C$  is an extremal ternary self-dual code of length  $n \equiv 0 \pmod{12}$  then the supports of all codewords of any nonzero weight  $w < n$  are the blocks of a 5-design.

## Theorem (Assmus and Mattson 1969)

The ternary extended quadratic residue codes  $QR^*$  of length  $n = 12, 24, 48$  and  $60$  are extremal and support 5-designs.

## Note

The code  $QR_{11}^*$  is equivalent to the extended ternary Golay code  $G_{12}$ .

# Pless Symmetry Codes

## Theorem (V. Pless 1969)

Let  $q \equiv -1 \pmod{3}$  be an odd prime power, and let

$$S_p = \begin{pmatrix} 0 & 1 & 1 & \cdots & 1 & \cdots & 1 \\ \chi(-1) & \chi(0) & \chi(1) & \cdots & \chi(\beta) & \cdots & \chi(-1) \\ \cdots & & & & & & \\ \chi(-1) & & & \cdots & \chi(\beta - \alpha) & \cdots & \\ \cdots & & & & & & \\ \chi(-1) & & & & & & \end{pmatrix},$$

where  $\chi(0) = 0$ ,  $\chi(a) = 1$  if  $a \neq 0$  is a square in  $GF(q)$ , and  $\chi(a) = -1$  if  $a \neq 0$  is not a square in  $GF(q)$ .

- The ternary code  $C(q)$  generated by  $(I_{q+1}, S_q)$  is self-dual.
- The codes  $C(q)$  for  $q = 5, 11, 17, 23, 29$  ( $n = 12, 24, 36, 48, 60$ ) are extremal and support 5-designs.

**Note.** The symmetry code  $C(5)$  is equivalent to the Golay code  $G_{12}$ .

# The known extremal ternary self-dual codes of length $n \equiv 0 \pmod{12}$

- $n = 12$ :  $G_{12} = QR_{11}^* = C(5)$ .
- $n = 24$ :  $QR_{23}^*$ ,  $C(11)$ .
- $n = 36$ :  $C(17)$ .
- $n = 48$ :  $QR_{47}^*$ ,  $C(23)$ .
- $n = 60$ :  $QR_{59}^*$ ,  $C(29)$ ,  $NV$ .

The code  $NV$  was found by G. Nebe and D. Villar in 2013 as a group theoretic analogue of the Pless symmetry code  $C(29)$ .

## Theorem

Up to equivalence, there is only one extremal ternary self-dual code of length 12 ( $G_{12}$ ), and two inequivalent codes of length 24 ( $QR_{23}^*$  and  $C(11)$ ).

# An analogue of the Pless symmetry codes

GABRIELE NEBE

nebe@math.rwth-aachen.de

DARWIN VILLAR

darwin.villar@rwth-aachen.de

Lehrstuhl D für Mathematik, RWTH Aachen University  
52056 Aachen, Germany

## Dedicated to the memory of Professor Stefan Dodunekov

**Abstract.** A series of monomial representations of  $SL_2(p)$  is used to construct a new series of self-dual ternary codes of length  $2(p+1)$  for all primes  $p \equiv 5 \pmod{8}$ . In particular we find a new extremal self-dual ternary code of length 60.

## 1 Introduction

In 1969 Vera Pless [6] discovered a family of self-dual ternary codes  $\mathcal{P}(p)$  of length  $2(p+1)$  for primes  $p$  with  $p \equiv -1 \pmod{6}$ . Together with the extended quadratic residue codes  $XQR(q)$  of length  $q+1$  ( $q$  prime,  $q \equiv \pm 1 \pmod{12}$ ) they define a series of self-dual ternary codes of high minimum distance (see [3, Chapter 16, §8]). For  $p=5$ , the Pless code  $\mathcal{P}(5)$  coincides with the Golay code  $\mathfrak{g}_{12}$  which is also the extended quadratic residue code  $XQR(11)$  of length 12.

Using invariant theory of finite groups, A. Gleason [2] has shown that the minimum distance of a self-dual ternary code of length  $4n$  cannot exceed  $3\lfloor \frac{n}{12} \rfloor + 3$ . Self-dual codes that achieve equality are called *extremal*. Both constructions, the Pless symmetry codes and the extended quadratic residue codes yield extremal ternary self-dual codes for small values of  $p$ .

This short note gives an interpretation of the Pless codes using monomial representations of the group  $SL_2(p)$ . This construction allows to read off a large subgroup of the automorphism group of the Pless codes (which was already described in [6]). A different but related series of monomial representations of  $SL_2(p)$  is investigated to construct a new series of self-dual ternary codes  $\mathcal{V}(p)$  of length  $2(p+1)$  for all primes  $p \equiv 5 \pmod{8}$ . The automorphism group of  $\mathcal{V}(p)$  contains the group  $SL_2(p)$ . For  $p=5$  we again find  $\mathcal{V}(5) \cong \mathfrak{g}_{12}$  the Golay code of length 12, but for larger primes these codes are new. In particular the code  $\mathcal{V}(29)$  is an extremal ternary code of length 60, so we now know three extremal ternary codes of length 60:  $XQR(59)$ ,  $\mathcal{P}(29)$  and  $\mathcal{V}(29)$ .

## 2 Codes and monomial groups

Let  $K$  be a field,  $n \in \mathbb{N}$ . Then the **monomial group**  $\text{Mon}_n(K^*) \leq GL_n(K)$  is the group of monomial  $n \times n$ -matrices over  $K$ , where a matrix is called

# Classification of ternary self-dual codes

- The largest length  $n \equiv 0 \pmod{4}$  for which **all** ternary self-dual codes have been classified up to equivalence, is  $n = 24$  (Harada and Munemasa 2009).
- The largest length  $n \equiv 0 \pmod{4}$  for which all **extremal** ternary self-dual codes have been classified up to equivalence, is  $n = 28$  (Harada, Munemasa and Venkov 2009).
- A partial classification of **extremal** ternary self-dual codes of length  $n \leq 40$  admitting automorphisms of **prime** order  $p \geq 5$  was given by C. W. Huffman (1992).
- G. Nebe (2012) proved that, up to equivalence, the only **extremal** ternary self-dual codes of length 48 that admit an automorphism of a **prime** order  $p \geq 5$ , are the Pless symmetry code and the extended QR code.
- Extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$  **do not exist** for  $n = 72, 96, 120$ , and all  $n \geq 144$ , because the weight enumerator contains a negative coefficient.

# Hadamard matrices and designs

A **Hadamard matrix** of order  $n$  is an  $n \times n$  matrix  $H$  of 1's and  $-1$ 's such that  $HH^T = nI$ , where  $I$  is the identity matrix.

It follows that  $n = 1, 2$ , or  $n = 4t$  for some integer  $t \geq 1$ .

An **automorphism** of a Hadamard matrix  $H$  is a pair of  $\{0, 1, -1\}$ -monomial matrices  $L, R$  such that  $LHR = H$ .

Two Hadamard matrices  $H_1, H_2$  of the same order are **equivalent** if there are monomial matrices  $L, R$  such that  $LH_1R = H_2$ .

A Hadamard matrix  $H$  is **normalized** with respect to its  $i$ th row and  $j$ th column if all entries in row  $i$  and column  $j$  are equal to 1.

If  $H$  is a Hadamard matrix of order  $n = 4t$  normalized with respect to row  $i$  and column  $j$ , deleting the  $i$ th row and the  $j$ th column and replacing all  $-1$ 's with zeros gives the  $(0, 1)$ -incidence matrix of a symmetric  $2$ - $(4t - 1, 2t - 1, t - 1)$  design  $D$  called a **Hadamard 2-design**.

If  $H$  is a Hadamard matrix of order  $n = 4t$  normalized with respect to row  $i$  and column  $j$ , deleting the  $j$ th column of  $H$  and the  $j$ th column of  $-H$  from the matrix  $(H, -H)$  gives the point-by-block  $(\pm 1)$ -incidence matrix of a Hadamard 3- $(4t, 2t, t - 1)$  design  $D^*$ .

A Hadamard matrix  $H$  of order  $n = 4t$  is **regular** if all rows of  $H$  contain the same number  $k$  of  $-1$ 's.

Then  $t = m^2$  for some integer  $m$ ,  $k = 2m^2 \pm m$ , and replacing all  $-1$ 's with zeros gives the  $(0, 1)$ -incidence matrix of a symmetric 2- $(4m^2, 2m^2 \pm t, m^2 \pm m)$  design (called sometimes a **Menon** design).

# Symmetry codes and Hadamard matrices

Let  $q$  be an odd prime power such that  $q \equiv -1 \pmod{3}$ .

## Theorem (Pless 1972)

- if  $q \equiv 1 \pmod{4}$ , the matrix

$$H_1(q) = \begin{pmatrix} I + S_q & -I + S_q \\ -I + S_q & -I - S_q \end{pmatrix} \quad (1)$$

is a Hadamard matrix whose rows are codewords from  $C(q)$ .

- If  $q \equiv 3 \pmod{4}$ , the matrix

$$H_3(q) = \begin{pmatrix} I + S_q & I + S_q \\ I - S_q & -I + S_q \end{pmatrix} \quad (2)$$

is a Hadamard matrix whose rows are codewords from  $C(q)$ .

The Hadamard matrices (1), (2) are equivalent to **Paley-Hadamard matrices of type II**.



# The automorphism group of $C(q)$

## Theorem (Pless 1972)

The symmetry code  $C(q)$  is invariant under a group of order  $\mathbf{q(q^2 - 1)}$  isomorphic to  $PGL(2, q)$ .

If  $q = 5, 11,$  or  $23$ , the rows of the Hadamard matrix  $H_1(q)$  from (1) (resp.  $H_3(q)$  from (2)) and  $-H_1(q)$  (resp.  $-H_3(q)$ ) exhaust all codewords of full weight, and span the code.

## Theorem (Pless 1972)

If  $q = 5, 11,$  or  $23$ , the full monomial automorphism group of  $C(q)$  coincides with the full automorphism group of  $H_1(q)$  (resp.  $H_3(q)$ ).

**Note.** The full automorphism group of a Paley-Hadamard matrix of type II for  $q > 5$  was determined by de Launey and Stafford in 2008, and is of order  $\mathbf{4fq(q^2 - 1)}$  if  $q = p^f$ , where  $p$  is prime.

**Note.** The symmetry code  $C(17)$  of length 36 contains 888 codewords of weight 36, while the number of codewords of weight 60 in  $C(29)$  is 41184.

## The code $L(q)$ : a monomial equivalent of $C(q)$

The sum of all rows of the generator matrix of the symmetry code  $C(q)$  is a vector  $v$  of full Hamming weight  $2q + 2$ , with all components equal to 1 if  $-1$  is not a square in  $GF(q)$ , and  $v$  has  $2q + 1$  components equal to 1, and the position labeled by  $\infty$  is equal to  $-1$  whenever  $-1$  is a square in  $GF(q)$ .

Next, we consider a code  $L(q)$  which is monomially equivalent to the Pless symmetry code  $C(q)$  and always contains the all-one vector, namely the code with a generator matrix  $G'$  given by

$$G' = (I_{q+1}, U_q), \quad (3)$$

where  $U_q$  is a  $(q + 1) \times (q + 1)$  matrix obtained from  $S_q$  by replacing every nonzero entry in the first column with  $-1$ . A parity check matrix of  $L(q)$  is given by

$$P = (-U_q^T, I_{q+1}). \quad (4)$$

# Symmetry codes and Hadamard 3-designs

## Theorem

The matrix  $H$  given by

$$H = \begin{pmatrix} G + P \\ G - P \end{pmatrix} = \begin{pmatrix} I_{q+1} - U_q^T & U_q + I_{q+1} \\ I_{q+1} + U_q^T & U_q - I_{q+1} \end{pmatrix} \quad (5)$$

is a Hadamard matrix with rows being full weight codewords of  $L(q)$ .

## Theorem

- The code  $L(q)$  contains a set of  $4q + 2$   $(0,1)$ -codewords of weight  $q + 1$  that form the block-by-point incidence matrix of a Hadamard  $3-(2q + 2, q + 1, (q - 1)/2)$  design  $D(q)$  associated with a Paley-Hadamard matrix of type II.
- If  $q = 5, 11, 17, 23$ , the code  $L(q)$  contains **exactly**  $4q + 2$   $(0,1)$ -codewords of weight  $q + 1$ , and every such codeword is the incidence vector of a block of the Hadamard 3-design  $D(q)$ .

# The permutation automorphism group of $L(q)$

## Theorem

If  $q = 5, 11, 17, 23$ , the full **permutation** automorphism group of  $L(q)$  coincides with the full automorphism group of the Hadamard 3-design  $D(q)$ , being of order  $q(q - 1)$ .

**Note.** The code  $L(29)$  contains 19606  $(0,1)$ -codewords of weight 30. It is an open question whether this set contains the incidence matrices of any Hadamard 3- $(60, 30, 14)$  designs that are not isomorphic to  $D(29)$ .

**Note.** The number of codewords of weight 60 in  $L(29)$  is 41184. It seems likely that there may be Hadamard matrices of order 60 formed by codewords of weight 60 that are not equivalent to the Paley-Hadamard matrix of type II.

## The code $L(17)$

The set of all 888 codewords of  $L(17)$  of full weight 36 comprises of the following disjoint subsets:

- the 36 rows of the Hadamard matrix  $H$  (5) normalized with respect to a row  $\bar{1}$ ;
- the 36 rows of  $2H$  that include a constant row  $\bar{2}$ ;
- a set  $T$  of 408 codewords having 15 components equal to 1 and 21 components equal to 2;
- a set  $2T$  of 408 codewords obtained by multiplying every codeword from  $T$  by 2.

**Note.** Adding  $\bar{2}$  to any  $(0, 1)$ -codeword of weight 18 gives a codeword of weight 36 with 18 1's and 18 2's; hence the code  $L(17)$  contains exactly 70  $(0, 1)$ -codewords of weight 18 obtained by adding the codeword  $\bar{2}$  to the rows of  $H$  and  $2H$ , and these 70  $(0, 1)$ -codewords form the incidence matrix of the 3-design  $D(17)$ .

# Hadamard matrices of order 36 contained in $L(17)$

## Theorem

- 1 The code  $L(17)$  contains **two** equivalence classes of Hadamard matrices of order 36:
  - a Hadamard matrix  $\mathbf{H}$  equivalent to a Paley-Hadamard matrix of type II, with full automorphism group of order  $19584 = 2^7 3^2 17$ ;
  - a second Hadamard matrix  $\mathbf{H}'$ , being a **regular** Hadamard matrix such that the associated **symmetric 2-** $(36, 15, 6)$  **design**  $D'$  has a **trivial** automorphism group.
- 2 The ternary code spanned by the incidence matrix of the  $2-$  $(36, 15, 6)$  design  $D'$  is an extremal ternary  $[36, 18, 12]$  code equivalent to the symmetry code  $C(17)$ .
- 3 The automorphism group of  $L(17)$  partitions the set of codewords of weight 36 into two orbits of length 72 and 816 respectively, the orbit of length 72 consisting of rows of  $H$  and  $-H$ .
- 4 The full automorphism group of the code  $L(17)$  coincides with the full automorphism group  $H$ .

# Some Hadamard matrices of order 36 and their codes

- Up to equivalence, there are exactly 11 Hadamard matrices of order 36 with automorphism groups of order divisible by 17 (Tonchev 1986). Each of these matrices spans a ternary self-dual code of length 36, but only one, namely the Paley-Hadamard matrix of type II, spans an extremal code, equivalent to the Pless symmetry code  $C(17)$ .
- Up to equivalence, there exists exactly one Hadamard matrix of order 36 with a doubly-transitive automorphism group, isomorphic to  $SP(6, 2) \times Z_2$  (N. Ito and J. Leon, 1980). This matrix spans a ternary self-orthogonal code of minimum distance 12 and dimension 14.

# Hadamard matrices and ternary QR codes

## Theorem.

- The symmetry codes  $C(11)$ ,  $C(23)$ , and  $C(29)$  have siblings with the same parameters and weight distribution, being ternary extended quadratic-residue codes.
- If  $q \equiv 3 \pmod{4}$  is a prime power, a quadratic residue (QR) code of length  $q$  is a code spanned by the  $(0,1)$ -incidence matrix  $A$  of a symmetric Hadamard  $2-(q, (q-1)/2, (q-3)/4)$  design associated with a Paley-Hadamard matrix of type I.
- The extended code is spanned by a matrix obtained by bordering  $A$  with the all-one column.
- If, in addition,  $q \equiv -1 \pmod{3}$ , that is,  $q = 12s + 11$ , the ternary extended QR code is self-dual.



## Theorem

Let  $q = 12s + 11$  be a prime power, and let  $QR_q$  be the ternary extended QR code of length  $q + 1$ .

- 1  $QR_q$  contains a **Paley-Hadamard matrix of type I** having as rows codewords of weight  $q + 1$ .
- 2  $QR_q$  contains a set of  $2q$   $(0,1)$ -codewords of weight  $(q + 1)/2$  that form the incidence matrix of a **Hadamard 3-design** associated with the Paley-Hadamard matrix of type I of order  $q + 1$ .
- 3 If  $q = 11, 23$  or  $47$ ,  $QR_q$  contains **exactly**  $2q$   $(0,1)$ -codewords of weight  $(q + 1)/2$ , and the permutation automorphism group of the code coincides with the full automorphism group of the Hadamard 3-design from part (2).

**Note.** The number of codewords of full weight 60 in  $QR_{59}$  is 41184. It is an interesting open question whether there are any Hadamard matrices of order 60 formed by codewords of weight 60 that are not equivalent to the Paley-Hadamard matrix of type I.

# Hadamard designs and self-dual codes

Hadamard matrices and designs have been used for the construction of self-orthogonal and self-dual codes over other finite fields.

- The extended binary Golay code is generated by a bordered incidence matrix of a symmetric Hadamard 2-(23, 11, 5) design associated with a Paley-Hadamard matrix of type I.
- Hadamard matrices of order 28 with an automorphism of order 7 were used for the classification of self-orthogonal codes over  $GF(7)$ :






V. Pless and V. D. Tonchev, Self-dual codes over  $GF(7)$ , *IEEE Trans. Info. Theory*, **33** (1987) 723-727.



V. D. Tonchev, Hadamard matrices of order 28 with an automorphism of order 7, *J. Combin. Theory Set A* **40** (1985) 62-81.

# Binary extremal self-dual codes derived from Hadamard matrices and designs

- The Paley-Hadamard matrix of type II of order 28 is the only Hadamard matrix of this order that admits an automorphism of order 13 and yields an extremal binary self-dual code of length 56:
  -  V. D. Tonchev, Hadamard matrices of order 28 with an automorphism of order 13, *J. Combin. Theory Set A* **35** (1983), 43-57.
  -  V. D. Tonchev, Self-orthogonal designs and extremal doubly-even codes, *J. Combin. Theory Set A* **52** (1989), 197-205.
- Many more extremal doubly-even binary self-dual codes derived from Hadamard matrices of order 28 were found in
  -  F. C. Bussemaker and V. D. Tonchev, New extremal doubly-even codes of length 56 derived from Hadamard matrices of order 28, *Discrete Math.* **76** (1989), 45-49.

# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?



# Open Questions

A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .

Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

# Open Questions








A general open question is whether there are more, yet unknown, extremal ternary self-dual codes of length  $n \equiv 0 \pmod{12}$ .








Specifically,

- 1 Is the symmetry code  $C(17)$  the only extremal code of length 36?
- 2 Are  $QR_{47}^*$  and  $C(23)$  the only extremal codes of length 48?
- 3 Are  $QR_{59}^*$ ,  $C(29)$ , and  $NV$  the only extremal codes of length 60?
- 4 Is there a Hadamard matrix of non-Paley type that spans  $QR_{59}^*$  or  $C(29)$ ?
- 5 Is the Nebe-Villar code the row space of a Hadamard matrix of order 60?
- 6 Is there an extremal ternary self-dual code of length 84, 108, or 132?

Thank You!

# Bibliography

-  E. F. Assmus, Jr. and H. F. Mattson, Jr., New 5-designs, *J. Combin. Theory, Ser. A* **6** (1969), 122-151.
-  W. Cary Huffman, On extremal self-dual ternary codes of lengths 28 to 40, *IEEE Trans. Info. Theory* **38** No. 4 (1992), 1395-1400.
-  W. de Launey, R. M. Stafford, On the automorphisms of Paley's type II Hadamard matrix, *Discrete Math.* **308** (2008), 2910-2924.
-  C. L. Mallows and N. J. A. Sloane, An upper bound for self-dual codes, *Information and Control* **22** (1973), 188-200.
-  G. Nebe, D. Villar, An analogue of the Pless symmetry codes, *Seventh International Workshop on Optimal Codes and Related Topics*, September 6 - 12, 2013, Albena, Bulgaria, pp. 158-163.
-  R. E. A. C. Paley, On orthogonal matrices, *J. Math. Phys.* **12** (1933) 311-320.
-  V. Pless, On a new family of symmetry codes and related new five-designs, *Bull. Amer. Math. Soc.* **75**, No. 6 (1969), 1339-1342

-  V. Pless, Symmetry codes over  $GF(3)$  and new five-designs, *J. Combin. Theory, Ser. A* **12** (1972), 119-142.
-  V. Pless and V. D. Tonchev, Self-dual codes over  $GF(7)$ , *IEEE Trans. Info. Theory*, **33** (1987) 723-727.
-  V. D. Tonchev, On Pless symmetry codes, ternary QR codes, and related Hadamard matrices and designs, *Designs, Codes and Cryprography*, to appear.
-  V. D. Tonchev, Hadamard matrices of order 28 with an automorphism of order 13, *J. Combin. Theory Set A* **35** (1983), 43-57.
-  V. D. Tonchev, Hadamard matrices of order 28 with an automorphism of order 7, *J. Combin. Theory Set A* **40** (1985) 62-81.
-  V. D. Tonchev, Self-orthogonal designs and extremal doubly-even codes, *J. Combin. Theory Set A* **52** (1989), 197-205.
-  V. D. Tonchev, Codes and Designs, Chapter 15 in: *Handbook of Coding Theory*, Vol. II, pages 1229-1268, V. S. Pless and W. C.